

5.3 Recursive Definitions

Σ finite set

Σ^* set of strings over Σ

(1) $\lambda \in \Sigma^*$ (λ denotes empty string)

(2) $w \in \Sigma^*$ and $x \in \Sigma$ implies $wx \in \Sigma^*$

Example $\Sigma = \{a, b\}$

$\Sigma^* = \{\lambda, a, b, aa, ab, ba, bb, aaa, aab, \dots\}$

The length of a string $w \in \Sigma^*$ is denoted $l(w)$ and defined by

$$l(\lambda) = 0$$

$$l(wx) = l(w) + 1 \quad \text{if } w \in \Sigma^*, x \in \Sigma$$

Example $l(a) = 1$ $l(a) = l(\lambda a) = l(\lambda) + 1$

$$l(ab) = 2 \quad l(ab) = l(a) + 1$$

$$= l(\lambda) + 1 + 1$$

$$= 2$$

For $x, y \in \Sigma^*$ let xy denote concatenation

$$x = aab, y = aba$$

$$xy = aababa$$

We will now consider structural induction

(1) Base case: Show the result holds for all base cases of recursive definition

(2) Recursive step: Show the result is true if it is true for each element used in recursion.

Example: Prove that $l(xy) = l(x) + l(y)$
for any $x, y \in \Sigma^*$

PF) Take any $x \in \Sigma^*$.

(1) We will look at $y = \epsilon$.

$$\text{LHS } l(xy) = l(x\epsilon) = l(x)$$

$$\text{RHS } l(x) + l(\epsilon) = l(x) + 0 = l(x)$$

(2) Assume that $l(xy) = l(x) + l(y)$ for some $y \in \Sigma^*$. Consider some $a \in \Sigma$.

$$\text{LHS } l(xya) = l(xy) + 1 = l(x) + l(y) + 1$$

$$\text{RHS } l(x) + l(ya) = l(x) + l(y) + 1$$

Ackerman's functions

Input $(m, n) \in \mathbb{N} \times \mathbb{N}$

$$A(m, n) = \begin{cases} 2n & m=0 \\ 0 & m \geq 1, n=0 \\ 2 & m \geq 1, n=1 \\ A(m-1, A(m, n-1)) & m \geq 1, n \geq 2 \end{cases}$$

$$A(0, 0) = 0$$

$$A(1, 0) = 0$$

$$A(0, 1) = 2$$

$$A(1, 1) = 2$$

$$A(1, 2) = A(0, A(1, 1)) = A(0, 2) = 4$$

$$\begin{aligned} \cancel{A(2, 2)} &= \cancel{A(1, A(2, 1))} = \cancel{A(1, 2)} = \cancel{A(0, A(0, 2))} \\ &= A \end{aligned}$$

$$A(2, 2) = A(1, A(2, 1))$$

$$A(2, 1) = 2$$

$$= A(1, 2)$$

$$= 4$$

Prove $A(m, 2) = 4 \quad \forall m \geq 1$

Prove $A(1, n) = 2^n \quad \forall n \geq 1$

Pf) Assume $A(m, 2) = 4$ for some $m \geq 1$. Notice $A(1, 2) = 4$ is the base case which we have checked.

$$\begin{aligned} A(m+1, 2) &= A(m, A(m+1, 1)) \\ &= A(m, 2) \\ &= 4 \end{aligned}$$

For a base case we see

$$A(1, 1) = 2 = 2^1$$

Now assume $A(1, n) = 2^n$ ~~then~~ for $n \geq 1$,

$$A(1, n+1) = A(0, A(1, n))$$

$$\begin{array}{|l} n \geq 1 \\ \hline n+1 \geq 2 \end{array}$$

$$= A(0, 2^n)$$

$$= 2 \cdot 2^n$$

$$= 2^{n+1}$$